## **Introduction to Petri Nets**

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## Welcome Guys :D

- Part 1: Intuitive Understanding of Petri nets
- Part 2: Basic Concepts
- Part 3: Elementary system nets
- Part 4: Sth interesting if we have time

## **Part 1: Intuitive Understanding**



Transition

# 4 Components





## **Part 1: Intuitive Understanding**



## **Part 1: Intuitive Understanding**



## Part 1: Intuitive Understanding Marking = State









#### Part 1: Intuitive Understan ding







## **Vincent's Automated Counter Design**



#### My Alternative design to this issue (without counter)











- Place: passive component
  - Store
  - Accumulate
  - Show

## Not A STATE

- **Transition**: active component
  - Consume
  - Transport
  - Produce
  - Change

- Arc
  - Bipartite Graph
  - "Either an arc runs from a place to a transition or the other way around" (Wolfgang, 14)





Net Structure

•

- N = (P, T, F)
  - P: set of all places
  - T: set of all transitions
  - $F \subseteq (P \times T) \cup (T \times P)$





•  $F \subseteq (P \times T) \cup (T \times P)$ 

- Pre-set and Post-set
  - In an unambiguously defined Net structure N, for a component x (place/transition) we can define
    - Pre-set of x: • $x =_{def} \{y \mid yFx\}$
    - Post-set of x:  $x^{\bullet} =_{def} \{y \mid xFy\}.$
    - Loop:  $x \in {}^{\bullet}y \text{ and } y \in {}^{\bullet}x$



• Marking



- Note:
  - All places must be considered.
  - Marking can be represented graphically. But not necessarily.

Initial marking

Distribute Cookie

- Multiset
  - Example: mixed kinds of tokens in a place.
    - $[\oplus \oplus \oplus \oplus \oplus]$
  - A multiset a is formally a mapping
  - a:U —> N

- Domain: Universe
  - Sufficiently Large
  - Collection of all examined tokens
- Codomain: Natural Number

- Multiset Example:
  - $U = \{ \bigoplus \ \bigcirc \ 0 \ 1 \ 2 \ 3 \ 4 \ 5 \ \}$
  - R = [ ⊕ ⊕ ⊕ ⊡ ⊡ ⊡ •]
    - a(•••) = 2
    - a(⊕) = 3
    - a(●) = 1

- a(u) = 0 for any other u in U
  - a(3) = 0
  - a(2) = 0

- $\mathcal{M}(U)$ : a set of all multisets over U
  - When U can be unambiguously identified, we write M(U)
  - Otherwise, we just write  $\mathcal{M}$ .

What about elementary system nets?
What is U?
What is M(U)?
Assume we allow at most one black
dot in one place,
do we still need the data structure multiset?

- Example: U = {
- 𝟸(U) = {[ ], [⊕], [⊡], [⊕⊕], [⊕⊕], [⊕⊕]
- A multiset a is finite if
  - $a(u) \neq 0$  for only finite number of  $u \in U$ .
  - +, -,  $\geq$ , $\leq$ ,=

- Marking: with the help of multisets:
   M: P -> M(U)
  - M<sub>0</sub>(coin slot) = [
    - a( • ) = 1
  - M<sub>0</sub>(cash box) = M<sub>0</sub>(signal) = M<sub>0</sub>(compartment)
    - a()=0

- M₀(storage) = [ ⊕ ⊕ ⊕ ⊕ ⊕ ]
  - a(⊕) = 5
- M<sub>0</sub>(counter) = [5]



- System Net
  - Net structure + Initial Marking + Transition condition (labelling in transition) + Arc\_labellings + cold\_transition
- Reachability of Marking: reachability is not so different from reachability of states in process graph especially in part 3.

• Final Marking

## **QUESTION TIME**



How many markings are there from the initial marking?

## **QUESTION TIME**

What about now?



## **QUESTION TIME**

What about now?









- Part1, Part2: Generic System Nets
- Part3: Elementary System Nets

- 2 Difference
  - Abstract Black Dot Token ONLY
  - No labelling: default is











$$M'(p) = \begin{cases} M(p) - 1 & \text{if } p \in {}^{\bullet}t \text{ and } p \notin t^{\bullet} \\ M(p) + 1 & \text{if } p \in t^{\bullet} \text{ and } p \notin {}^{\bullet}t \\ M(p) & \text{otherwise} \end{cases}$$

Step Rule

$$M_0 \xrightarrow{t_1} M_1 \xrightarrow{t_2} \ldots \xrightarrow{t_n} M_n$$

(Wolfgang, 26)











# More...

- Labelling (Generic System Nets)
- Petri Net modelling ME
- Code and Hot Transition
- Petri Net Modelling Relay Race
- Petri Net modelling Vivian\_reading\_books\_unless\_rain
- Some interesting design in cookie vending machine

#### Labelling of Arcs

• "Represent the tokens that flow through the arc <u>at the occurrence</u> <u>of transition." (Wolfgang, 16)</u>

 $\overline{pt}$  or  $\overline{tp}$ 

- <u>NOT REAL TOKENs</u>
  - Constant
  - Variable/Function



- Labelling of Transitions
  - Condition with variable.
  - Evaluated to True or False



Models ME using elementary system nets. Find sth wrong here?





Relay Race: Process Graph



Relay Race: Elementary System Nets



Relay Race: Generic System Nets





#### Models Vivian's Read\_book\_unless\_rainy days



## Links to recordings



### References



Reisig, Wolfgang. Understanding Petri Nets: Modeling Techniques, Analysis Methods, Case Studies. Berlin, Heidelberg: Springer Berlin Heidelberg: Imprint: Springer, 2013. Print.

# Thank you!

